

# Pointwise Qualitative Factorization of 2-forms into 1-forms over $\mathbb{R}^4$

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Abstract: Assume an arbitrary 2-form over  $\mathbb{R}^4$ , and a fixed point  $x$  in  $\mathbb{R}^4$  for which the signs  $[+,0,-]$  of all six coordinates of the 2-form are known, are given. If the 2-form factors into a wedge product of two 1-forms, a complete classification of the possible combinations of signs of the coordinates of the 1-forms (at  $x$ ) is presented. Also, it is shown that some sign combinations in the coordinates of the 2-form preclude the 2-form being factorizable into a wedge product of two 1-forms.

Why would we want this problem solved?

An elementary example is relativistic gyroscopic precession. A spinning body (an electron, the Earth, etc.) that accelerates because forces act on its center of mass precesses its axis of rotation according to the formula

$$d\mathbf{S}/d\tau = (\mathbf{u} \wedge \mathbf{a}) \cdot \mathbf{S}$$

where  $\tau$  is proper time measured along the object's worldline,  $\mathbf{S}$  is the angular-momentum 4-vector of the object,  $\mathbf{u}$  is the 4-velocity of the object, and  $\mathbf{a}$  is the 4-acceleration applied to the center of mass of the object. The pointwise algebraic constraints discussed in this paper are constraints on the 2-form, which is known to be a wedge product of two 1-forms. [From another point of view, the angular momentum is Fermi-Walker transported.]

By definition, all 4-vectors, 1-forms, and 2-forms mentioned are over a 4-dimensional space-time manifold, which locally cannot be distinguished from  $\mathbb{R}^4$ . Thus, the pointwise algebraic constraints, to be discussed, are applicable in any space-time manifold, except at singularities.

### Problem definition

Use the standard basis for  $\mathbb{R}^4$ , and the induced standard basis for 1-forms and 2-forms over  $\mathbb{R}^4$ . Denote the coordinates (in  $\mathbb{R}$ ) of a fixed 2-form  $\omega$  (at a fixed point  $x$  in  $\mathbb{R}^4$ ), by

$$\omega(x) = c_{12} dx^1 \wedge dx^2 + c_{13} dx^1 \wedge dx^3 + c_{14} dx^1 \wedge dx^4 + c_{23} dx^2 \wedge dx^3 + c_{24} dx^2 \wedge dx^4 + c_{34} dx^3 \wedge dx^4$$

Also, (assuming that 1-forms  $\alpha, \beta$  exist such that  $\omega = \alpha \wedge \beta$ ), denote these in coordinates (in  $\mathbb{R}$ ) by

$$\alpha(x) = a_1 dx^1 + a_2 dx^2 + a_3 dx^3 + a_4 dx^4$$

$$\beta(x) = b_1 dx^1 + b_2 dx^2 + b_3 dx^3 + b_4 dx^4$$

Recall that in (assumed possible) the coordinate computation of  $\omega = \alpha \wedge \beta$ ,

$$c_{12} = a_1 b_2 - a_2 b_1 \quad c_{13} = a_1 b_3 - a_3 b_1 \quad c_{14} = a_1 b_4 - a_4 b_1$$

$$c_{34} = a_3 b_4 - a_4 b_3 \quad c_{24} = a_2 b_4 - a_4 b_2 \quad c_{23} = a_2 b_3 - a_3 b_2$$

We sometimes use  $i, j = 1..4$ . [ $i < j$ ] when used as subscripts in  $c_{ij}$ ; otherwise,  $i$  and  $j$  are not necessarily related. Note that in the above translation, using the permutation  $[ab]$  on the above equations is the coordinate version of writing  $-\omega = \beta \wedge \alpha$ . We will use this freely when using the computational lemmas [i.e.: the permutation  $[ab]$  may be applied globally in a lemma, provided all of the  $c_{ij}$  are negated.]

Finally, we require that it be known, for each of the 2-form coordinates  $c_{ij}$ , whether it is positive, zero, or negative [ $c_{ij} > 0$ ,  $c_{ij} = 0$ ,  $c_{ij} < 0$ ]. Parts of the classification will consider two of these cases at once, for various  $c_{ij}$ .

From this information, we want to know:

- Is it possible to rule out a factorization  $\omega = \alpha \wedge \beta$ , just from the sign information for the  $c_{ij}$ ?
- If we have failed to rule out such a factorization  $\omega = \alpha \wedge \beta$ , what does the sign information for the  $c_{ij}$  allow us to infer about sign information for the  $a_i, b_j$ ?

### Computing the action of the symmetric group $S_4$ on the coordinates

It will be useful, in reducing the amount of work in the classification, to have an explicit description of the action of the symmetric group  $S_4$  on the 2-form coordinates  $c_{ij}$ , as well as the effects of this action on the constraints. [For the rest of this paper,  $S_4$  shall always refer to the symmetric group  $S_4$ .]

The overall results are as follows:

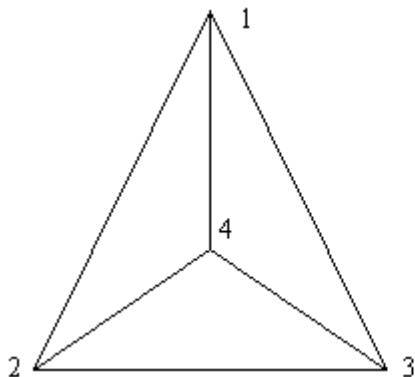
id	[12]	[13]	[14]	[23]	[24]	[34]	[123]	[132]	[124]	[142]	[134]
$c_{12}$	$-c_{12}$	$-c_{23}$	$-c_{24}$	$c_{13}$	$c_{14}$	$c_{12}$	$c_{23}$	$-c_{13}$	$c_{24}$	$-c_{14}$	$-c_{23}$
$c_{13}$	$c_{23}$	$-c_{13}$	$-c_{34}$	$c_{12}$	$c_{13}$	$c_{14}$	$-c_{12}$	$-c_{23}$	$c_{23}$	$-c_{34}$	$c_{34}$
$c_{14}$	$c_{24}$	$c_{34}$	$-c_{14}$	$c_{14}$	$c_{12}$	$c_{13}$	$c_{24}$	$c_{34}$	$-c_{12}$	$-c_{24}$	$-c_{13}$
$c_{23}$	$c_{13}$	$-c_{12}$	$c_{23}$	$-c_{23}$	$-c_{34}$	$c_{24}$	$-c_{13}$	$c_{12}$	$-c_{34}$	$c_{14}$	$c_{24}$
$c_{24}$	$c_{14}$	$c_{24}$	$-c_{12}$	$c_{34}$	$-c_{24}$	$c_{23}$	$c_{34}$	$c_{14}$	$-c_{14}$	$c_{12}$	$-c_{12}$
$c_{34}$	$c_{34}$	$c_{14}$	$-c_{13}$	$c_{24}$	$-c_{23}$	$-c_{34}$	$c_{14}$	$c_{24}$	$-c_{13}$	$-c_{23}$	$-c_{14}$

[143]	[234]	[243]	[12] [34]	[13] [24]	[14] [23]	[1234]	[1243]	[1324]	[1342]	[1423]	[1432]
$-c_{24}$	$c_{13}$	$c_{14}$	$-c_{12}$	$c_{34}$	$-c_{34}$	$c_{23}$	$c_{24}$	$c_{34}$	$-c_{13}$	$-c_{34}$	$-c_{14}$
$-c_{14}$	$c_{14}$	$c_{12}$	$c_{24}$	$-c_{13}$	$-c_{24}$	$c_{24}$	$-c_{12}$	$-c_{23}$	$c_{34}$	$-c_{14}$	$-c_{24}$
$-c_{34}$	$c_{12}$	$c_{13}$	$c_{23}$	$-c_{23}$	$-c_{14}$	$-c_{12}$	$c_{23}$	$-c_{13}$	$-c_{23}$	$-c_{24}$	$c_{34}$
$-c_{12}$	$c_{34}$	$-c_{24}$	$c_{14}$	$-c_{14}$	$-c_{23}$	$c_{34}$	$-c_{14}$	$-c_{24}$	$c_{14}$	$-c_{13}$	$c_{12}$
$c_{23}$	$-c_{23}$	$-c_{34}$	$c_{13}$	$-c_{24}$	$-c_{13}$	$-c_{13}$	$-c_{34}$	$-c_{14}$	$c_{12}$	$-c_{23}$	$c_{13}$
$c_{13}$	$-c_{24}$	$c_{23}$	$-c_{34}$	$c_{12}$	$-c_{12}$	$c_{14}$	$c_{13}$	$-c_{12}$	$-c_{24}$	$c_{12}$	$c_{23}$

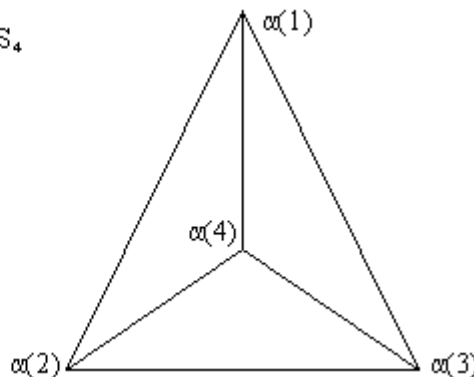
A graphical method of computing the above action is illustrated as follows:

The edges with vertices  $i,j$  index the coefficient  $c_{ij}$ , in standard  $i$  orientation. It will be convenient to define [for  $1 \leq i < j \leq 4$ ]

$$c_{ji} := -c_{ij}$$

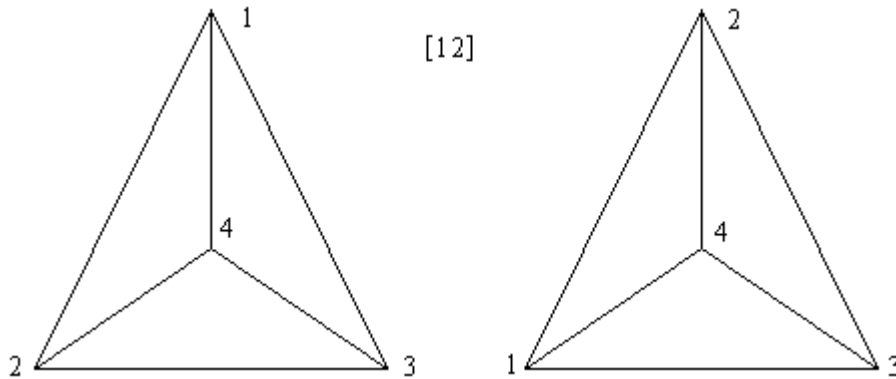


$\alpha \in S_4$



Then we can write that  $\alpha \in S_4$  maps  $c_{ij}$  to  $c_{\alpha(i)\alpha(j)}$ .

As an example, consider the image of the transposition [12]:



Either the graphical method, or the algebraic-table method, may be used in calculating the equivalence classes of the various cases in the classification.

### Computational facts

Besides the sign information about  $a_i b_j$  relative to  $a_i$  and  $b_j$ , and  $c_{ij}$  relative to  $a_i b_j - a_j b_i$ ,

	$b_j < 0$	$b_j = 0$	$b_j > 0$			$a_j b_i < 0$	$a_j b_i = 0$	$a_j b_i > 0$
$a_i < 0$	$a_i b_j > 0$	$a_i b_j = 0$	$a_i b_j < 0$		$a_i b_j < 0$	$c_{ij} ? 0$	$c_{ij} < 0$	$c_{ij} < 0$
$a_i = 0$	$a_i b_j = 0$	$a_i b_j = 0$	$a_i b_j = 0$		$a_i b_j = 0$	$c_{ij} > 0$	$c_{ij} = 0$	$c_{ij} < 0$
$a_i > 0$	$a_i b_j < 0$	$a_i b_j = 0$	$a_i b_j > 0$		$a_i b_j > 0$	$c_{ij} > 0$	$c_{ij} > 0$	$c_{ij} ? 0$

Note:  $c_{ij} < 0$  implies

- [ $a_i < 0$  AND  $b_j > 0$ ]
- OR [ $a_i > 0$  AND  $b_j < 0$ ]
- OR [ $a_j < 0$  AND  $b_i < 0$ ]
- OR [ $a_j > 0$  AND  $b_i > 0$ ]

Note:  $c_{ij} > 0$  implies

- [ $a_i < 0$  AND  $b_j < 0$ ]
- OR [ $a_i > 0$  AND  $b_j > 0$ ]
- OR [ $a_j < 0$  AND  $b_i > 0$ ]
- OR [ $a_j > 0$  AND  $b_i < 0$ ]

Note:  $c_{ij}=0$  implies

- [ $a_i < 0$  AND  $a_j < 0$  AND  $b_i < 0$  AND  $b_j < 0$ ]
- OR [ $a_i > 0$  AND  $a_j < 0$  AND  $b_i < 0$  AND  $b_j > 0$ ]
- OR [ $a_i < 0$  AND  $a_j > 0$  AND  $b_i > 0$  AND  $b_j < 0$ ]
- OR [ $a_i > 0$  AND  $a_j > 0$  AND  $b_i > 0$  AND  $b_j > 0$ ]
- OR [ $a_i = 0$  AND  $a_j = 0$ ]
- OR [ $a_i = 0$  AND  $b_i = 0$ ]
- OR [ $b_j = 0$  AND  $a_j = 0$ ]
- OR [ $b_j = 0$  AND  $b_i = 0$ ]
- OR [ $a_i < 0$  AND  $a_j < 0$  AND  $b_i > 0$  AND  $b_j > 0$ ]
- OR [ $a_i > 0$  AND  $a_j < 0$  AND  $b_i > 0$  AND  $b_j < 0$ ]
- OR [ $a_i < 0$  AND  $a_j > 0$  AND  $b_i < 0$  AND  $b_j > 0$ ]
- OR [ $a_i > 0$  AND  $a_j > 0$  AND  $b_i < 0$  AND  $b_j < 0$ ]

Note [there is an analog for  $a_i \neq 0 \neq a_j$ ]:

$$c_{ij} = a_i b_j - a_j b_i > 0 \text{ iff } \frac{a_i}{b_i} > \frac{a_j}{b_j}$$

$$b_i \neq 0 \neq b_j \text{ IMPLIES } [c_{ij} = a_i b_j - a_j b_i = 0 \text{ iff } \frac{a_i}{b_i} = \frac{a_j}{b_j}]$$

$$c_{ij} = a_i b_j - a_j b_i < 0 \text{ iff } \frac{a_i}{b_i} < \frac{a_j}{b_j}$$

Note: the configuration  $c_{ij}=0=c_{ik}$ ,  $c_{jk} \neq 0$ , implies that at least one of  $a_i$ ,  $a_j$ ,  $a_k$  is zero, and also that at least one of  $b_i$ ,  $b_j$ , and  $b_k$  is zero. [Otherwise, the calculational note about comparing quotients applies, and  $c_{jk}=0$  — contradiction.]. As a shorthand, we refer to instances of the above comment as ALO1(i,j,k); the arguments do not commute, and  $j < k$  by convention.

In fact, the conditions yielding ALO1(i,j,k) actually imply that  $[a_i=0=b_i]$ . [This will be Lemma 2, later.] We will prove this directly at Lemma 2. A way to visualize this is that the cross-product of two vectors in  $\mathbb{R}^3$  is the dual of the wedge product of the vectors (reinterpreted coordinate-wise as 1-forms) over  $\mathbb{R}^3$ . The hypothesis then translates to requiring the cross-product to be a vector that is non-zero only in the first coordinate; since the cross-product, geometrically, is perpendicular to both of the vectors it was computed from, both of the vectors it was computed from have first coordinate zero. [Actually, we will (re)prove that the visualization is a correct interpretation of the cross-product over  $\mathbb{R}^3$ .]

## Computational Lemmas

Lemma 0

$$\begin{aligned}
 & [[a_i=0 \text{ OR } b_j=0] \text{ AND } c_{ij}<0] \text{ IMPLIES } a_j b_i > 0 \\
 & [[a_i=0 \text{ OR } b_j=0] \text{ AND } c_{ij}=0] \text{ IMPLIES } a_j b_i = 0 \\
 & [[a_i=0 \text{ OR } b_j=0] \text{ AND } c_{ij}>0] \text{ IMPLIES } a_j b_i < 0 \\
 & [[a_j=0 \text{ OR } b_i=0] \text{ AND } c_{ij}<0] \text{ IMPLIES } a_i b_j < 0 \\
 & [[a_j=0 \text{ OR } b_i=0] \text{ AND } c_{ij}=0] \text{ IMPLIES } a_i b_j = 0 \\
 & [[a_j=0 \text{ OR } b_i=0] \text{ AND } c_{ij}>0] \text{ IMPLIES } a_i b_j > 0
 \end{aligned}$$

Sketch of calculational proof: element chase of the first two tables in the computational facts. The first three use the hypothesis to single out the second row of the RHS table, while the last three use the hypothesis to single out the second column of the RHS table.

Lemma 1

$$[[a_i=0 \text{ OR } [b_j=0 \text{ AND } b_k=0]] \text{ AND } c_{ij} \neq 0 \text{ AND } c_{ik}=0] \text{ IMPLIES } a_k=0$$

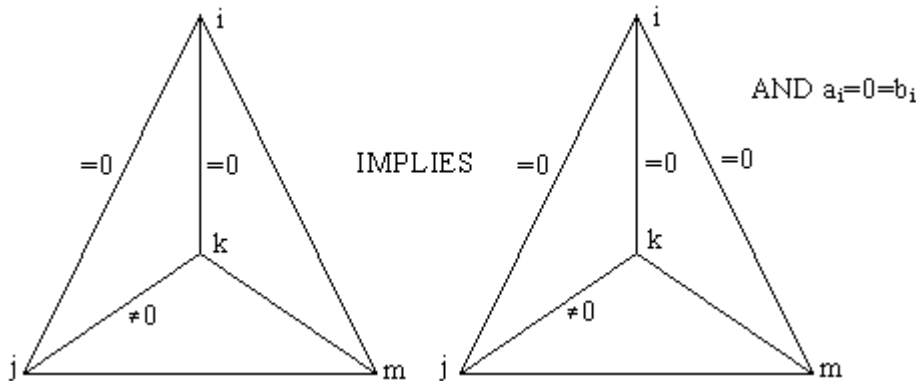
Sketch of calculational proof:

Recall that  $c_{ij}=a_i b_j - a_j b_i$ . Use Lemma 0 with  $c_{ij} \neq 0$  to get  $a_j b_i \neq 0$  [making  $b_i \neq 0$ ], and with  $c_{ik}=0$  to get  $a_k b_i = 0$ . Since  $b_i \neq 0$ , and  $R$  has no zero divisors,  $a_k = 0$ .

[NOTE: other variations on the above hypotheses [two 2-form coordinates with a common index, and a variable entering into at least one of the coordinate calculations, are constrained] lead to repeats of instances of lemma 0. I have suppressed the repeats of lemma 0 in lemma 1.]

Lemma 2:

$$[c_{ij}=0=c_{ik} \text{ AND } c_{jk} \neq 0] \text{ IMPLIES } [a_i=0=b_i \text{ AND } c_{im}=0]$$



Sketch of calculational proof:

We have (using calculational notes about the hypothesized entries):

[  
    [ $a_i < 0$  AND  $a_j < 0$  AND  $b_i < 0$  AND  $b_j < 0$ ]  
    OR [  $a_i > 0$  AND  $a_j < 0$  AND  $b_i < 0$  AND  $b_j > 0$  ]  
    OR [  $a_i < 0$  AND  $a_j > 0$  AND  $b_i > 0$  AND  $b_j < 0$  ]  
    OR [  $a_i > 0$  AND  $a_j > 0$  AND  $b_i > 0$  AND  $b_j > 0$  ]  
    OR [  $a_i = 0$  AND  $a_j = 0$  ]  
    OR [  $a_i = 0$  AND  $b_i = 0$  ]  
    OR [  $b_j = 0$  AND  $a_j = 0$  ]  
    OR [  $b_j = 0$  AND  $b_i = 0$  ]  
    OR [  $a_i < 0$  AND  $a_j < 0$  AND  $b_i > 0$  AND  $b_j > 0$  ]  
    OR [  $a_i > 0$  AND  $a_j < 0$  AND  $b_i > 0$  AND  $b_j < 0$  ]  
    OR [  $a_i < 0$  AND  $a_j > 0$  AND  $b_i < 0$  AND  $b_j > 0$  ]  
    OR [  $a_i > 0$  AND  $a_j > 0$  AND  $b_i < 0$  AND  $b_j < 0$  ] ]  
AND [  
    [  $a_i < 0$  AND  $a_k < 0$  AND  $b_i < 0$  AND  $b_k < 0$  ]  
    OR [  $a_i > 0$  AND  $a_k < 0$  AND  $b_i < 0$  AND  $b_k > 0$  ]  
    OR [  $a_i < 0$  AND  $a_k > 0$  AND  $b_i > 0$  AND  $b_k < 0$  ]  
    OR [  $a_i > 0$  AND  $a_k > 0$  AND  $b_i > 0$  AND  $b_k > 0$  ]  
    OR [  $a_i = 0$  AND  $a_k = 0$  ]  
    OR [  $a_i = 0$  AND  $b_i = 0$  ]  
    OR [  $b_k = 0$  AND  $a_k = 0$  ]  
    OR [  $b_k = 0$  AND  $b_i = 0$  ]  
    OR [  $a_i < 0$  AND  $a_k < 0$  AND  $b_i > 0$  AND  $b_k > 0$  ]  
    OR [  $a_i > 0$  AND  $a_k < 0$  AND  $b_i > 0$  AND  $b_k < 0$  ]  
    OR [  $a_i < 0$  AND  $a_k > 0$  AND  $b_i < 0$  AND  $b_k > 0$  ]  
    OR [  $a_i > 0$  AND  $a_k > 0$  AND  $b_i < 0$  AND  $b_k < 0$  ] ]  
AND [  
    [  $a_j \neq 0$  AND  $b_k \neq 0$  ]  
    OR [  $a_k \neq 0$  AND  $b_j \neq 0$  ] ]

The third OR-clause implies (but not conversely) that  $[a_j \neq 0 \text{ OR } b_j \neq 0]$  and  $[a_k \neq 0 \text{ OR } b_k \neq 0]$ .

Applying this to the first and second OR-clauses yields:

$$\begin{aligned}
 & [ \\
 & \quad \text{OR} \quad [a_i < 0 \text{ AND } a_j < 0 \text{ AND } b_i < 0 \text{ AND } b_j < 0] \\
 & \quad \text{OR} \quad [a_i > 0 \text{ AND } a_j < 0 \text{ AND } b_i < 0 \text{ AND } b_j > 0] \\
 & \quad \text{OR} \quad [a_i < 0 \text{ AND } a_j > 0 \text{ AND } b_i > 0 \text{ AND } b_j < 0] \\
 & \quad \text{OR} \quad [a_i > 0 \text{ AND } a_j > 0 \text{ AND } b_i > 0 \text{ AND } b_j > 0] \\
 & \quad \text{OR} \quad [a_i = 0 \text{ AND } a_j = 0] \\
 & \quad \text{OR} \quad [a_i = 0 \text{ AND } b_i = 0] \\
 & \quad \text{OR} \quad [b_j = 0 \text{ AND } b_i = 0] \\
 & \quad \text{OR} \quad [a_i < 0 \text{ AND } a_j < 0 \text{ AND } b_i > 0 \text{ AND } b_j > 0] \\
 & \quad \text{OR} \quad [a_i > 0 \text{ AND } a_j < 0 \text{ AND } b_i > 0 \text{ AND } b_j < 0] \\
 & \quad \text{OR} \quad [a_i < 0 \text{ AND } a_j > 0 \text{ AND } b_i < 0 \text{ AND } b_j > 0] \\
 & \quad \text{OR} \quad [a_i > 0 \text{ AND } a_j > 0 \text{ AND } b_i < 0 \text{ AND } b_j < 0] \quad ] \\
 \text{AND} & [ \\
 & \quad \text{OR} \quad [a_i < 0 \text{ AND } a_k < 0 \text{ AND } b_i < 0 \text{ AND } b_k < 0] \\
 & \quad \text{OR} \quad [a_i > 0 \text{ AND } a_k < 0 \text{ AND } b_i < 0 \text{ AND } b_k > 0] \\
 & \quad \text{OR} \quad [a_i < 0 \text{ AND } a_k > 0 \text{ AND } b_i > 0 \text{ AND } b_k < 0] \\
 & \quad \text{OR} \quad [a_i > 0 \text{ AND } a_k > 0 \text{ AND } b_i > 0 \text{ AND } b_k > 0] \\
 & \quad \text{OR} \quad [a_i = 0 \text{ AND } a_k = 0] \\
 & \quad \text{OR} \quad [a_i = 0 \text{ AND } b_i = 0] \\
 & \quad \text{OR} \quad [b_k = 0 \text{ AND } b_i = 0] \\
 & \quad \text{OR} \quad [a_i < 0 \text{ AND } a_k < 0 \text{ AND } b_i > 0 \text{ AND } b_k > 0] \\
 & \quad \text{OR} \quad [a_i > 0 \text{ AND } a_k < 0 \text{ AND } b_i > 0 \text{ AND } b_k < 0] \\
 & \quad \text{OR} \quad [a_i < 0 \text{ AND } a_k > 0 \text{ AND } b_i < 0 \text{ AND } b_k > 0] \\
 & \quad \text{OR} \quad [a_i > 0 \text{ AND } a_k > 0 \text{ AND } b_i < 0 \text{ AND } b_k < 0] \quad ] \\
 \text{AND} & [ \\
 & \quad \text{OR} \quad [a_j \neq 0 \text{ AND } b_k \neq 0] \\
 & \quad \text{OR} \quad [a_k \neq 0 \text{ AND } b_j \neq 0] \quad ]
 \end{aligned}$$

Now,  $\text{ALO1}(i,j,k)$  implies that when we expand the AND of the first two OR-clauses: the first four, and the last four, of the terms in the first OR-clause cannot successfully expand with any of the eleven terms in the second OR-clause. [The first four and the last four terms of the second OR-clause violate  $\text{ALO1}(i,j,k)$ , and the central three terms of the second OR-clause directly contradict the term from the first OR-clause that we are expanding.] Similar reasoning holds with the first four, and the last four, terms of the second OR-clause with respect to the first OR-clause.

This allows reducing the above to:

$$\begin{aligned}
 & [ \\
 & \quad \text{OR} \quad [a_i = 0 \text{ AND } a_j = 0] \\
 & \quad \text{OR} \quad [a_i = 0 \text{ AND } b_i = 0] \\
 & \quad \text{OR} \quad [b_j = 0 \text{ AND } b_i = 0] \quad ] \\
 \text{AND} & [ \\
 & \quad \text{OR} \quad [a_i = 0 \text{ AND } a_k = 0] \\
 & \quad \text{OR} \quad [a_i = 0 \text{ AND } b_i = 0] \\
 & \quad \text{OR} \quad [b_k = 0 \text{ AND } b_i = 0] \quad ] \\
 \text{AND} & [ \\
 & \quad \text{OR} \quad [a_j \neq 0 \text{ AND } b_k \neq 0] \\
 & \quad \text{OR} \quad [a_k \neq 0 \text{ AND } b_j \neq 0] \quad ]
 \end{aligned}$$



Now, expand the AND of the first two OR-clauses:

$$\begin{aligned}
 & [ \quad \quad \quad [[a_i=0 \text{ AND } a_j=0] \text{ AND } [a_i=0 \text{ AND } a_k=0]] \\
 & \quad \text{OR} \quad [[a_i=0 \text{ AND } a_j=0] \text{ AND } [a_i=0 \text{ AND } b_i=0]] \\
 & \quad \text{OR} \quad [[a_i=0 \text{ AND } a_j=0] \text{ AND } [b_k=0 \text{ AND } b_i=0]] \\
 & \quad \text{OR} \quad [[a_i=0 \text{ AND } b_i=0] \text{ AND } [a_i=0 \text{ AND } a_k=0]] \\
 & \quad \text{OR} \quad [[a_i=0 \text{ AND } b_i=0] \text{ AND } [a_i=0 \text{ AND } b_i=0]] \\
 & \quad \text{OR} \quad [[a_i=0 \text{ AND } b_i=0] \text{ AND } [b_k=0 \text{ AND } b_i=0]] \\
 & \quad \text{OR} \quad [[b_j=0 \text{ AND } b_i=0] \text{ AND } [a_i=0 \text{ AND } a_k=0]] \\
 & \quad \text{OR} \quad [[b_j=0 \text{ AND } b_i=0] \text{ AND } [a_i=0 \text{ AND } b_i=0]] \\
 & \quad \text{OR} \quad [[b_j=0 \text{ AND } b_i=0] \text{ AND } [b_k=0 \text{ AND } b_i=0]] \quad ] \\
 \text{AND} & [ \quad \quad \quad [a_j \neq 0 \text{ AND } b_k \neq 0] \\
 & \quad \text{OR} \quad [a_k \neq 0 \text{ AND } b_j \neq 0] \quad ]
 \end{aligned}$$

i.e. (Redundancy elimination only, no sorting, remove second duplicate term: first OR-clause)

$$\begin{aligned}
 & [ \quad \quad \quad [a_i=0 \text{ AND } a_j=0 \text{ AND } a_k=0] \\
 & \quad \text{OR} \quad [a_i=0 \text{ AND } a_j=0 \text{ AND } b_i=0] \\
 & \quad \text{OR} \quad [a_i=0 \text{ AND } a_j=0 \text{ AND } b_k=0 \text{ AND } b_i=0] \\
 & \quad \text{OR} \quad [a_i=0 \text{ AND } b_i=0 \text{ AND } a_k=0] \\
 & \quad \text{OR} \quad [a_i=0 \text{ AND } b_i=0] \\
 & \quad \text{OR} \quad [a_i=0 \text{ AND } b_i=0 \text{ AND } b_k=0] \\
 & \quad \text{OR} \quad [b_j=0 \text{ AND } b_i=0 \text{ AND } a_i=0 \text{ AND } a_k=0] \\
 & \quad \text{OR} \quad [b_j=0 \text{ AND } b_i=0 \text{ AND } a_i=0] \\
 & \quad \text{OR} \quad [b_j=0 \text{ AND } b_i=0 \text{ AND } b_k=0] \quad ] \\
 \text{AND} & [ \quad \quad \quad [a_j \neq 0 \text{ AND } b_k \neq 0] \\
 & \quad \text{OR} \quad [a_k \neq 0 \text{ AND } b_j \neq 0] \quad ]
 \end{aligned}$$

Now, the second OR-clause also yields  $[a_j \neq 0 \text{ OR } a_k \neq 0]$  and  $[b_j \neq 0 \text{ OR } b_k \neq 0]$ . This removes the first and last clauses of the first OR-clause:

$$\begin{aligned}
 & [ \quad \quad \quad [a_i=0 \text{ AND } a_j=0 \text{ AND } b_i=0] \\
 & \quad \text{OR} \quad [a_i=0 \text{ AND } a_j=0 \text{ AND } b_k=0 \text{ AND } b_i=0] \\
 & \quad \text{OR} \quad [a_i=0 \text{ AND } b_i=0 \text{ AND } a_k=0] \\
 & \quad \text{OR} \quad [a_i=0 \text{ AND } b_i=0] \\
 & \quad \text{OR} \quad [a_i=0 \text{ AND } b_i=0 \text{ AND } b_k=0] \\
 & \quad \text{OR} \quad [b_j=0 \text{ AND } b_i=0 \text{ AND } a_i=0 \text{ AND } a_k=0] \\
 & \quad \text{OR} \quad [b_j=0 \text{ AND } b_i=0 \text{ AND } a_i=0] \quad ] \\
 \text{AND} & [ \quad \quad \quad [a_j \neq 0 \text{ AND } b_k \neq 0] \\
 & \quad \text{OR} \quad [a_k \neq 0 \text{ AND } b_j \neq 0] \quad ]
 \end{aligned}$$

We can now factor  $[a_i=0 \text{ AND } b_i=0]$  out of the first OR-clause:

$$\begin{aligned}
 & [a_i=0 \text{ AND } b_i=0] \\
 \text{AND } [ & \quad [a_j=0] \\
 & \quad \text{OR } [a_j=0 \text{ AND } b_k=0] \\
 & \quad \text{OR } [a_k=0] \\
 & \quad \text{OR } [\text{true}] \\
 & \quad \text{OR } [b_k=0] \\
 & \quad \text{OR } [b_j=0 \text{ AND } a_k=0] \\
 & \quad \text{OR } [b_j=0] ] \\
 \text{AND } [ & \quad [a_j \neq 0 \text{ AND } b_k \neq 0] \\
 & \quad \text{OR } [a_k \neq 0 \text{ AND } b_j \neq 0] ]
 \end{aligned}$$

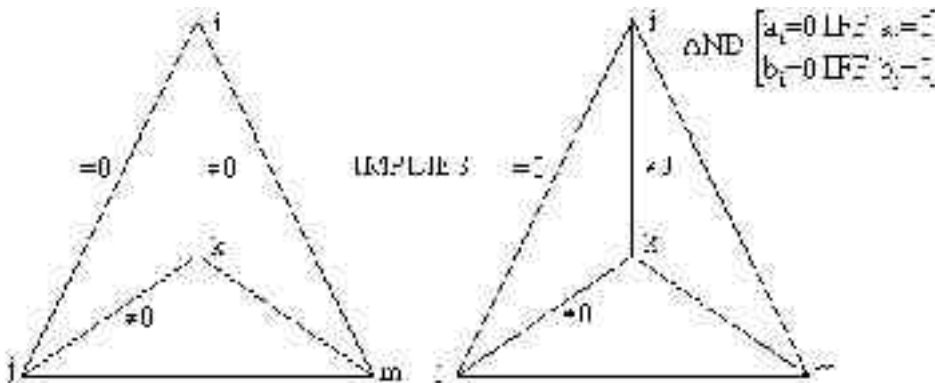
i.e. [the first OR-clause is trivially true]

$$\begin{aligned}
 & [a_i=0 \text{ AND } b_i=0] \\
 \text{AND } [ & \quad [a_j \neq 0 \text{ AND } b_k \neq 0] \\
 & \quad \text{OR } [a_k \neq 0 \text{ AND } b_j \neq 0] ]
 \end{aligned}$$

It is now trivial that  $c_{im}=0$ .

Lemma 3:

$$\begin{aligned}
 & [c_{ij}=0 \text{ AND } c_{ik} \neq 0 \text{ AND } c_{jk} \neq 0] \text{ IMPLIES } [a_i=0 \text{ IFF } a_j=0] \\
 & \quad \text{swap } a, b \\
 & [c_{ij}=0 \text{ AND } c_{ik} \neq 0 \text{ AND } c_{jk} \neq 0] \text{ IMPLIES } [b_i=0 \text{ IFF } b_j=0]
 \end{aligned}$$



Sketch of calculational proof:

Assume  $c_{ij}=0, c_{ik} \neq 0 \neq c_{jk}$ . Then Lemma 1 yields both ' $a_i=0$  implies  $a_j=0$ ' [ $c_{ik} \neq 0, c_{ij}=0$ ] and ' $a_j=0$  implies  $a_i=0$ ' [ $c_{jk} \neq 0, c_{ij}=0$ ], i.e. ' $a_i=0$  iff  $a_j=0$ '.

Lemma 4:

$$c_{13}c_{24} = c_{12}c_{34} + c_{14}c_{23}$$

Sketch of calculational proof:

$$\begin{aligned} c_{12}c_{34}+c_{14}c_{23} &= [a_1b_2-a_2b_1][a_3b_4-a_4b_3]+[a_1b_4-a_4b_1][a_2b_3-a_3b_2] \\ &= a_1b_2a_3b_4 - a_1b_2a_4b_3 - a_2b_1a_3b_4 + a_2b_1a_4b_3 + a_1b_4a_2b_3 - a_1b_4a_3b_2 - a_4b_1a_2b_3 + a_4b_1a_3b_2 \\ &= -a_1b_2a_4b_3 - a_2b_1a_3b_4 + a_1b_4a_2b_3 + a_4b_1a_3b_2 = [a_1b_3 - a_3b_1][a_2b_4 - a_4b_2] = c_{13}c_{24} \end{aligned}$$

It is a routine check that this equation is fixed under the  $S_4$  vertex action. [The simplest way is to verify that the sign of all three terms is negated under the transpositions [14], [24], [34] — which are a generating set for  $S_4$ .]

### The classification

We will first do a preliminary classification, based on whether  $c_{ij} = 0$ , or  $c_{ij} \neq 0$ .

We first enumerate the distinct cases (for the  $c_{ij}$ ) under the action of  $S_4$  on the indexes 1..4, and then explicitly describe the possible cases for the  $a_i$  and  $b_j$ , for each case of the  $c_{ij}$ .

We first enumerate representatives of the equivalence classes of the constraints:

Partial case	$c_{12}$	$c_{13}$	$c_{14}$	$c_{23}$	$c_{24}$	$c_{34}$	Images under $S_4$ action	Total cases per image	Total cases indexed by partial case
0	$\neq 0$	$\neq 0$	$\neq 0$	$\neq 0$	$\neq 0$	$\neq 0$	1	$2^6=64$	64
1	$=0$	$\neq 0$	$\neq 0$	$\neq 0$	$\neq 0$	$\neq 0$	6	$2^5=32$	192
2A	$=0$	$=0$	$\neq 0$	$\neq 0$	$\neq 0$	$\neq 0$	12	$2^4=16$	192
2B	$=0$	$\neq 0$	$\neq 0$	$\neq 0$	$\neq 0$	$=0$	3	$2^4=16$	48
3A	$=0$	$=0$	$=0$	$\neq 0$	$\neq 0$	$\neq 0$	4	$2^3=8$	32
3B	$=0$	$\neq 0$	$\neq 0$	$=0$	$\neq 0$	$=0$	12	$2^3=8$	96
3C	$=0$	$=0$	$\neq 0$	$=0$	$\neq 0$	$\neq 0$	4	$2^3=8$	32
4A	$=0$	$\neq 0$	$=0$	$=0$	$\neq 0$	$=0$	3	$2^2=4$	12
4B	$=0$	$=0$	$=0$	$=0$	$\neq 0$	$\neq 0$	12	$2^2=4$	48
5	$=0$	$=0$	$=0$	$=0$	$=0$	$\neq 0$	6	$2^1=2$	12
6	$=0$	$=0$	$=0$	$=0$	$=0$	$=0$	1	$2^0=1$	1

The last column is for how many cases in the total classification are indexed by the above partial cases. The left-hand factor is the number of images of the representative under the  $S_4$  action on the vertices; the right-hand factor is 2 to the number of  $\neq 0$  constraints in the partial case.

As a cross-check, note that the sum of the counts is  $729 [= 3^6]$ , the number of cases to be dealt with in the total classification. Also, note that all of these cases are preserved by commuting the 1-forms in the wedge product [coordinate-wise, swapping a and b with the permutation [ab].] We

will use this symmetry freely, without comment, in the preliminary classification.

Note: Cases 2A, 3B, and 4A directly violate both Lemma 2 and Lemma 4, thus cannot factor into a wedge product of 1-forms. We only need to analyze the other cases.

**Case 0:**  $[c_{12} \neq 0, c_{13} \neq 0, c_{14} \neq 0, c_{23} \neq 0, c_{24} \neq 0, c_{34} \neq 0]$   
 [Symmetry on vertices:  $S_4$ ]

In Case 0, from all 6  $c_{ij} \neq 0$  we have:

$[a_1 \neq 0 \text{ AND } b_2 \neq 0] \text{ OR } [a_2 \neq 0 \text{ AND } b_1 \neq 0]$   
 $[a_1 \neq 0 \text{ AND } b_3 \neq 0] \text{ OR } [a_3 \neq 0 \text{ AND } b_1 \neq 0]$   
 $[a_1 \neq 0 \text{ AND } b_4 \neq 0] \text{ OR } [a_4 \neq 0 \text{ AND } b_1 \neq 0]$   
 $[a_2 \neq 0 \text{ AND } b_3 \neq 0] \text{ OR } [a_3 \neq 0 \text{ AND } b_2 \neq 0]$   
 $[a_2 \neq 0 \text{ AND } b_4 \neq 0] \text{ OR } [a_4 \neq 0 \text{ AND } b_2 \neq 0]$   
 $[a_3 \neq 0 \text{ AND } b_4 \neq 0] \text{ OR } [a_4 \neq 0 \text{ AND } b_3 \neq 0]$

Lemma 4 rules out the following 16 subcases (either verify the  $S_4$  vertex action and the kernel, or all 16 cases):

```

Computation of images under vertex S4 action of
c_12<0, c_13<0, c_14<0, c_23<0, c_24>0, c_34<0
      12 13 14 23 24 34
[id]      < < < < > < [234],[243]      // kernel
[23]      < < < > < > [24],[34]
[14]      < > > < > > [1423],[1432]
[13][24]  < > > > < < [132],[134]
[12]      > < > < < < [1243],[1234]
[14][23]  > < > > > > [142],[143]
[12][34]  > > < < < > [123],[124]
[13]      > > < > > < [1324],[1342]
  
```

This does **not** contain its sign-flip, qualitatively.

Thus, the commuted case is depicted by:

```

Dual computation of images under vertex S4 action of
c_12>0, c_13>0, c_14>0, c_23>0, c_24<0, c_34>0
      12 13 14 23 24 34
[id]      > > > > < > [234],[243]      // kernel
[23]      > > > < > < [24],[34]
[14]      > < < > < < [1423],[1432]
[13][24]  > < < < > > [132],[134]
[12]      < > < > > > [1243],[1234]
[14][23]  < > < < < < [142],[143]
[12][34]  < < > > > < [123],[124]
[13]      < < > < < > [1324],[1342]
  
```

**Case 1:**  $[c_{12}=0, c_{13} \neq 0, c_{14} \neq 0, c_{23} \neq 0, c_{24} \neq 0, c_{34} \neq 0]$   
 [Symmetry on vertices:  $\langle [12],[34] \rangle$  subgroup  $S_4$ ]

Lemma 3 yields  $[c_{12}=0, c_{13}\neq 0, c_{14}\neq 0]$ :

$$\begin{aligned} a_1=0 &\text{ IFF } a_2=0 \\ b_1=0 &\text{ IFF } b_2=0 \end{aligned}$$

From  $c_{13}\neq 0, c_{14}\neq 0, c_{23}\neq 0, c_{24}\neq 0, c_{34}\neq 0$ , we have (after using above):

$$\begin{aligned} [a_1\neq 0 \text{ AND } b_3\neq 0] &\text{ OR } [a_3\neq 0 \text{ AND } b_1\neq 0] \\ [a_1\neq 0 \text{ AND } b_4\neq 0] &\text{ OR } [a_4\neq 0 \text{ AND } b_1\neq 0] \\ [a_3\neq 0 \text{ AND } b_4\neq 0] &\text{ OR } [a_4\neq 0 \text{ AND } b_3\neq 0] \end{aligned}$$

Lemma 4 rules out the following subcases (again, either verify the  $S_4$  action and the kernels, or all of the cases):

---

	12	13	14	23	24	34		
...								
Type II: failed Lemma 4	=	<	<	<	>	<	[id]	// kernel
	=	<	>	<	<	<	[12]	
	=	<	<	>	<	>	[34]	
	=	>	<	<	<	>	[12] [34]	
Type III: failed Lemma 4	=	<	<	<	>	>	[id]	// kernel
	=	<	>	<	<	>	[12]	
	=	<	<	>	<	<	[34]	
	=	>	<	<	<	<	[12] [34]	
...								
Type III sign-flip: failed Lemma 4	=	>	>	>	<	<	[id]	// kernel
	=	>	<	>	>	<	[12]	
	=	>	>	<	>	>	[34]	
	=	<	>	>	>	>	[12] [34]	
Type II sign-flip: failed Lemma 4	=	>	>	>	<	>	[id]	// kernel
	=	>	<	>	>	>	[12]	
	=	>	>	<	>	<	[34]	
	=	<	>	>	>	<	[12] [34]	

---

**Case 2B:**  $[c_{12}=0, c_{13}\neq 0, c_{14}\neq 0, c_{23}\neq 0, c_{24}\neq 0, c_{34}=0]$

[Symmetry on vertices:  $\langle [12], [34], [13][24] \rangle$  subgroup  $S_4$ ]

Lemma 3 yields  $[c_{12}=0, c_{13}\neq 0, c_{14}\neq 0, c_{23}\neq 0, c_{24}\neq 0, c_{34}=0]$ :

$$\begin{aligned} a_1=0 &\text{ IFF } a_2=0 \\ b_1=0 &\text{ IFF } b_2=0 \\ a_3=0 &\text{ IFF } a_4=0 \\ b_3=0 &\text{ IFF } b_4=0 \end{aligned}$$

From  $c_{13}\neq 0$ , we have:

$$[a_1\neq 0 \text{ AND } b_3\neq 0] \text{ OR } [a_3\neq 0 \text{ AND } b_1\neq 0]$$

Lemma 4 rules out the following subcases (again, either verify the  $S_4$  action and the kernels, or all of the cases):

---

	12	13	14	23	24	34		
...								
Type II: failed Lemma 4	=	<	<	<	>	=	[id]	// kernel
	=	<	>	<	<	=	[12]	
	=	<	<	>	<	=	[34]	
	=	>	<	<	<	=	[12] [34]	
...								
Type II sign-flip: failed Lemma 4	=	>	>	>	<	=	[id]	// kernel
	=	>	<	>	>	=	[12]	
	=	>	>	<	>	=	[34]	
	=	<	>	>	>	=	[12] [34]	

---

**Case 3A:**  $[c_{12}=0, c_{13}=0, c_{14}=0, c_{23}\neq 0, c_{24}\neq 0, c_{34}\neq 0]$   
 [Symmetry on vertices:  $\text{Perm}(\{2,3,4\})$  subgroup  $S_4$ ]

Lemma 2  $[c_{12}=0, c_{13}=0, c_{23}\neq 0]$  yields for Case 3A:

$$a_1=0=b_1$$

We also have, from  $c_{23}\neq 0, c_{24}\neq 0, c_{34}\neq 0$ :

$$\begin{aligned} & [a_2\neq 0 \text{ AND } b_3\neq 0] \text{ OR } [a_3\neq 0 \text{ AND } b_2\neq 0] \\ & [a_2\neq 0 \text{ AND } b_4\neq 0] \text{ OR } [a_4\neq 0 \text{ AND } b_2\neq 0] \\ & [a_3\neq 0 \text{ AND } b_4\neq 0] \text{ OR } [a_4\neq 0 \text{ AND } b_3\neq 0] \end{aligned}$$

**Case 3C:**  $[c_{12}=0, c_{13}=0, c_{14}\neq 0, c_{23}=0, c_{24}\neq 0, c_{34}\neq 0]$   
 [Symmetry on vertices:  $S_3$  subgroup  $S_4$ ]

Lemma 3 yields  $[c_{12}=0, c_{14}\neq 0, c_{24}\neq 0; c_{13}=0, c_{14}\neq 0, c_{34}\neq 0; c_{23}=0, c_{24}\neq 0, c_{34}\neq 0]$ :

$$\begin{aligned} a_1=0 & \text{ IFF } a_2=0 \text{ IFF } a_3=0 \\ b_1=0 & \text{ IFF } b_2=0 \text{ IFF } b_3=0 \end{aligned}$$

From  $c_{14}\neq 0$ , we also have:

$$[a_1\neq 0 \text{ AND } b_4\neq 0] \text{ OR } [a_4\neq 0 \text{ AND } b_1\neq 0]$$

**Case 4B:**  $[c_{12}=0, c_{13}=0, c_{14}=0, c_{23}=0, c_{24}\neq 0, c_{34}\neq 0]$   
 [Symmetry on vertices:  $\langle [23] \rangle$  subgroup  $S_4$ ]

Lemma 2 yields  $[c_{12}=0, c_{14}=0, c_{24}\neq 0]$ :

$$a_1=0=b_1$$

Lemma 3 yields  $[c_{23}=0, c_{24}\neq 0, c_{34}\neq 0]$ :

$$a_2=0 \text{ IFF } a_3=0$$

$$b_2=0 \text{ IFF } b_3=0$$

Also, from  $c_{24} \neq 0$ , we have:

$$[a_2 \neq 0 \text{ AND } b_4 \neq 0] \text{ OR } [a_4 \neq 0 \text{ AND } b_2 \neq 0]$$

**Case 5:** [ $c_{12}=0, c_{13}=0, c_{14}=0, c_{23}=0, c_{24}=0, c_{34} \neq 0$ ]  
 [Symmetry on vertices:  $\langle [12], [34] \rangle$  subgroup  $S_4$ ]

Lemma 2 yields [ $c_{13}=0, c_{14}=0, c_{34} \neq 0; c_{23}=0, c_{24}=0, c_{34} \neq 0$ ]:  
 $a_1=0=b_1$   
 $a_2=0=b_2$

Also, from  $c_{34} \neq 0$ , we have:

$$[a_3 \neq 0 \text{ AND } b_4 \neq 0] \text{ OR } [a_4 \neq 0 \text{ AND } b_3 \neq 0]$$

**Case 6:** [ $c_{12}=0, c_{13}=0, c_{14}=0, c_{23}=0, c_{24}=0, c_{34}=0$ ]  
 [Symmetry on vertices:  $S_4$ ]

We have, for Case 6:

$$\alpha=0 \text{ OR } \beta=0$$

OR all of these:

$$a_1=0 \text{ IFF } b_1=0$$

$$a_2=0 \text{ IFF } b_2=0$$

$$a_3=0 \text{ IFF } b_3=0$$

$$a_4=0 \text{ IFF } b_4=0$$

Sketch of calculational proof:

Since  $0 \wedge \beta = 0 = \alpha \wedge 0$ , we need only consider the case when  $\alpha \neq 0 \neq \beta$ . In this case, by the problem symmetry we need only to prove  $a_1=0 \text{ IFF } b_1=0$ .

Suppose  $a_1=0$ . Then, by definition of  $c_{ij}$  [ $j=2..4$ ],  $a_j b_1=0$ . Since  $\alpha \neq 0$ , one of  $a_j \neq 0$ , yielding  $b_1=0$ . This proves  $a_1=0 \text{ IMPLIES } b_1=0$ ; swap a and b to get parallel reasoning for the other direction (yielding  $a_1=0 \text{ IFF } b_1=0$ ).

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